CSCI235 Database Systems

Normalization in Practice

Dr Janusz R. Getta

School of Computing and Information Technology -University of Wollongong

```
A relational schema R = (A, B, C)
```

```
Functional dependencies: AB \rightarrow C
```

Keys?

```
If AB \rightarrow C is valid in R and it covers entire relational schema then its left hand side is a minimal key (A, B)
```

```
No other minimal keys
```

Normal form?

```
Left hand side of AB \rightarrow C is a minimal key K = (A, B)
```

BCNF

A relational schema R =(A, B, C)

Functional dependencies: $AB \rightarrow C, C \rightarrow B$

Keys?

If $AB \rightarrow C$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A, B)

If $C \rightarrow B$ then through augmentation rule $AC \rightarrow AB$

If $AC \rightarrow AB$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A, C)

Normal form?

Not BCNF because left hand side of $C \rightarrow B$ is not a minimal key

3NF because right hand side of $C \rightarrow B$ is a prime attribute

Decomposition into **BCNF**?

R1 = (B, C), R2 = (A, C) or

R1 = (B, C), R2 = (A, B) Created by Janusz R. Getta, CSCI235 Database Systems, Autumn 2024

3/20

A relational schema R = (A, B, C)

```
Functional dependencies: AB \rightarrow C, C \rightarrow B, C \rightarrow A
```

Keys?

If $AB \rightarrow C$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A, B)

If $C \rightarrow B$ and $C \rightarrow A$ then through union rule $C \rightarrow AB$

If $C \rightarrow AB$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (C)

Normal form?

BCNF because left hand side of each functional dependency is a minimal key

```
A relational schema R = (A, B, C)
```

```
Functional dependencies: A \rightarrow B
```

Keys?

If $A \rightarrow B$ is valid in R then through augmentation rule $AC \rightarrow BC$

If $AC \rightarrow BC$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A, C)

Normal form?

not 2NF because a nonprime attribute B functionally depends (A \rightarrow B) on a subset of primary key (A, C)

Decomposition into **BCNF**?

R1 = (A, B), R2 = (A, C) or

```
R1 = (A, B), R2 = (B, C)
```

- A relation al schema R = (A, B, C)
- Functional dependencies: $A \rightarrow B, B \rightarrow A$

Keys?

If A \rightarrow B then through augmentation rule AC \rightarrow BC

If $AC \rightarrow BC$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A, C)

If $B \rightarrow A$ then through augmentation rule $BC \rightarrow AC$

If $BC \rightarrow AC$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (B, C)

Normal form?

Not BCNF because left hand side of $A \rightarrow B$ is not a minimal key

3NF because right hand side of $A \rightarrow B$ is a prime attribute and right hand side of $B \rightarrow A$ is a prime attribute

Decomposition into **BCNF**?

R1 = (A, B), R2 = (A, C) or

R1 = (A, B), R2 = (B, C)

```
A relational schema R = (A, B, C)
```

Functional dependencies: $A \rightarrow B, B \rightarrow C$

Keys?

If A \rightarrow B and B \rightarrow C then through transitivity rule A \rightarrow C

If A \rightarrow B and A \rightarrow C then through union rule A \rightarrow BC

If $A \rightarrow BC$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A)

Normal form?

Not 3NF because a non prime attribute C is transitively dependent on primary key A

2NF because no nonprime attribute depends on a part of primary key

Decomposition into **BCNF**?

- R1 = (A, B), R2 = (B, C) or
- R1 = (A, B), R2 = (A, C)



Created by Janusz R. Getta, CSCI235 Database Systems, Autumn 2024

- A relational schema R =(A, B, C, D)
- Functional dependencies: $A \rightarrow B, A \rightarrow C, B \rightarrow D$

Keys?

If $A \rightarrow B$ and $A \rightarrow C$ then through union rule $A \rightarrow BC$

If A \rightarrow B and B \rightarrow D then through transitivity rule A \rightarrow D

If A \rightarrow BC and A \rightarrow D then through union rule A \rightarrow BCD

If $A \rightarrow BCD$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A)

Normal form?

Not 3NF because a non prime attribute D is transitively dependent on primary key A

2NF because no nonprime attribute depends on a part of primary key

Decomposition into **BCNF**?

- R1 = (A, B, C), R2 = (B, D) or
- R1 = (A, B, C), R2 = (A, D)



Created by Janusz R. Getta, CSCI235 Database Systems, Autumn 2024

- A relational schema R =(A, B, C, D)
- Functional dependencies: $A \rightarrow B$, $B \rightarrow D$, $C \rightarrow B$

Keys?

- If A \rightarrow B and B \rightarrow D then through transitivity rule A \rightarrow D
- If A \rightarrow D and A \rightarrow B then through union rule A \rightarrow BD
- If A \rightarrow BD then through augmentation rule AC \rightarrow BCD

If $AC \rightarrow BCD$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A, C)

- If $C \rightarrow B$ and $B \rightarrow D$ then through transitivity rule $C \rightarrow D$
- If $C \rightarrow D$ and $C \rightarrow B$ then through union rule $C \rightarrow BD$
- If C \rightarrow BD then through augmentation rule AC \rightarrow ABD

If $AC \rightarrow BCD$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A, C)

Normal form?

Not 2NF because a nonprime attribute B depends on a part of a primary key (A, C)

Decomposition into **BCNF**?

R1 = (A, B), R2 = (B, C), R2 = (B, D)

- A relational schema R =(A, B, C, D)
- Functional dependencies: $A \rightarrow B$, $A \rightarrow C$, $B \rightarrow A$, $B \rightarrow C$

Keys?

- If A \rightarrow B and A \rightarrow C then through union rule A \rightarrow BC
- If A \rightarrow BC then through augmentation rule AD \rightarrow BCD

If $AD \rightarrow BCD$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A, D)

- If $B \rightarrow A$ and $B \rightarrow C$ then through union rule $B \rightarrow AC$
- If $B \rightarrow AC$ then through augmentation rule $BD \rightarrow ACD$

If $BD \rightarrow ACD$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (B, D)

Normal form?

Not 2NF because a nonprime attribute C depends on a part of a primary key (B, D)

Decomposition into **BCNF**?

R1 = (A, B), R2 = (B, C), R2 = (A, D)

- A relational schema R = (A, B, C, D)
- Functional dependencies: $AB \rightarrow C, C \rightarrow D, D \rightarrow A, D \rightarrow B$

Keys?

- If $AB \rightarrow C$ and $C \rightarrow D$ then through transitivity rule $AB \rightarrow D$
- If AB \rightarrow D and AB \rightarrow C then through union rule AB \rightarrow CD

If $AB \rightarrow CD$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A, B)

- If $D \rightarrow A$ and $D \rightarrow B$ then through union rule $D \rightarrow AB$
- If D \rightarrow AB and AB \rightarrow C then through transitivity rule D \rightarrow C
- If $D \to C$ and $D \to AB$ then $D \to ABC$

If $D \rightarrow ABC$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (D)

- If $C \rightarrow D$ and $D \rightarrow AB$ then through transitivity rule $C \rightarrow AB$
- If $C \rightarrow D$ and $C \rightarrow AB$ then through union rule $C \rightarrow ABD$

If $C \rightarrow ABD$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (C)

Normal form?

BCNF because left hand side of each functional dependency is a superky

- A relational schema R = (A, B, C, D)
- Functional dependencies: $A \rightarrow B$, $B \rightarrow C$, $C \rightarrow D$, $D \rightarrow A$

Keys?

- If A \rightarrow B and B \rightarrow C then through transitivity rule A \rightarrow C
- If A \rightarrow C and C \rightarrow D then through transitivity rule A \rightarrow D

If $A \rightarrow B$ and $A \rightarrow C$ and $A \rightarrow D$ then through union rule $A \rightarrow BCD$

If $A \rightarrow BCD$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A)

If $B \rightarrow C$ and $C \rightarrow D$ then through transitivity rule $B \rightarrow D$

If $B \rightarrow D$ and $D \rightarrow A$ then through transitivity rule $B \rightarrow A$

If $B \rightarrow C$ and $B \rightarrow D$ and $B \rightarrow A$ then through union rule $B \rightarrow ACD$

If $B \rightarrow ACD$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (B)

- If $C \rightarrow D$ and $D \rightarrow A$ then through transitivity rule $C \rightarrow A$
- If $C \rightarrow A$ and $A \rightarrow B$ then through transitivity rule $C \rightarrow B$
- If $C \rightarrow A$ and $C \rightarrow B$ and $C \rightarrow D$ then through union rule $C \rightarrow ABD$

If $C \rightarrow ABD$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (C)

If $D \rightarrow A$ and $A \rightarrow B$ then through transitivity rule $D \rightarrow B$

If $D \rightarrow B$ and $B \rightarrow C$ then through transitivity rule $D \rightarrow C$

If $D \rightarrow A$ and $D \rightarrow B$ and $D \rightarrow C$ then through union rule $D \rightarrow ABC$

If $D \rightarrow ABC$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (D)

Normal form ?

BCNF because left hand side of each functional dependency is a superky

References

T. Connoly, C. Begg, Database Systems, A Practical Approach to Design, Implementation, and Management, Chapter 14.5 The Process of Normalization, Chapter 14.6 First Normal Form (1NF), Chapter 14.7 Second Normal Form (2NF), Chapter 14.8 Third Normal Form (3NF), Chapter 14.9 General definitions of 2NF and 3NF, Chapter 15.2 Boyce-Codd Normal Form (BCNF), Chapter 15.3 Review of Normalization Up to BCNF, Pearson Education Ltd, 2015